

COT 6401 The Analysis of Algorithms
Test (March 11, 2009)
Due: **midnight of March 12, 2009**

Name _____ SSN _____

1. (20%) (**Greedy**) A telecom company needs to install base stations to cover all houses along a long road. These houses can be sparsely distributed along the road. Suppose the coverage of each base station is 5 miles. Design an *optimal solution* that covers all houses using as few base stations as possible. Prove that your algorithm is optimal.
2. (20%) (**Linear programming**) Solve the following linear program using SIMPLEX and show all the relevant steps:

$$\text{maximize } x_1 + 2x_2$$

subject to

$$4x_1 - x_2 \leq 9$$

$$x_1 + x_2 \leq 8$$

$$5x_1 - 2x_2 \geq -3$$

$$x_1, x_2 \geq 0$$

Provide a geometric explanation of the solution by plotting the corresponding *feasible region* in a 2-D space.

3. (20%) (**Divide-and-conquer**) Suppose the only way to access a database of student GPA is through a simple query k and that the system returns the k^{th} smallest value that it contains. Design an algorithm that finds the *median GPA* from two separate databases A (with m values) and B (with n values) using at most $\Theta(\log(m+n))$ queries. Show explicitly how your solution meets the requirement. Note that the median GPA is the $\lceil (m+n)/2 \rceil^{\text{th}}$ smallest value in A and B .
4. (20%) (**Brute-force**) Let $G = (V, E)$ be a k -nary tree with n nodes. The distance between two nodes in G is the length of the path connecting these two nodes (neighbors have distance 1). The *diameter* of G is the maximal distance over all pairs of nodes. Design a *linear-time solution* (i.e. $\Theta(n)$) to find the diameter of G .
5. (20%) (**Dynamic programming**) Design an *optimal solution* using dynamic programming for the *general coin changing problem*. Let a coin of denomination i , $1 \leq i \leq n$, have value d_i . Use the example with three coins with values 1, 4, and 6 units to illustrate the correctness of your solution by showing optimal results for changes from 1 to 10.
6. (**Bonus: 20%**) Quicksort can be modified to find the k^{th} *smallest element* from n elements so that in most cases it does much less work than is needed to sort the set completely.
 - (a) Write a modified quicksort algorithm for this purpose.
 - (b) Show that when this algorithm is used to find the median, the worst case is $\Theta(n^2)$.
 - (c) Develop a recurrence equation for the average running time of this algorithm.
 - (d) Analyze the average running time of the algorithm. What is the asymptotic order?